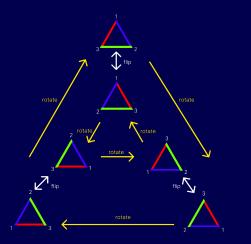
Finite Group Theory: the major problems, and why we care



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The formal definition of a group is one we've all seen many times:

Definition

A group is a pair (G, \circ) where G is a set, and $\circ : G \times G \to G$ is a binary operation on G that satisfies the following three properties:

- (i) Associativity. For all $g, h, k \in G$, $(g \circ h) \circ k = g \circ (h \circ k)$.
- (ii) **Identity**. G has an element, which we call 1_G , satisfying $g \circ 1_G = 1_G \circ g = g$ for all $g \in G$. We call 1_G an *identity element* in (G, \circ) .
- (iii) **Inverse**. For all $g \in G$, there exists an element in G, which we call g^{-1} , satisfying $g \circ g^{-1} = g^{-1} \circ g = 1_G$, where 1_G is as in (ii). We call g^{-1} an *inverse* of g in (G, \circ) .

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Most branches of algebra we study today (i.e. Linear Algebra; Ring Theory; Group Theory; Module Theory, etc.) were built from a desire to get rigorous answers to questions from other areas of mathematics and science.

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Let's look back at the definition:

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Part (i) represents the fact that, if I perform two symmetries (g and h) and then perform another (k) some time later; this is the same as performing g, waiting a while, then performing h and k.

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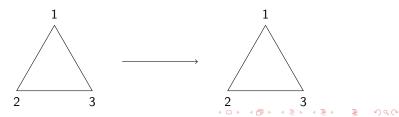
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In part (ii), the identity represents the "do nothing" symmetry of an object.



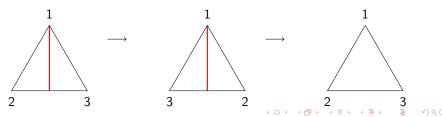
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Part (iii) encodes the fact the "reverse" of every symmetry is still a symmetry.



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More generally, the group of symmetries of an n element set (or equivalently the number of shuffles of a pack of n cards) is called the symmetric group of degree n, and is written S_n .

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Historical note: The first definition of a group was given by Galois in 1830, and it was less abstract than the one above. Indeed, Galois defined a group of substitutions of degree n to be what we now call a subgroup of the symmetric group S_n . This shows that the study of groups is fundamentally motivated by the desire to understand symmetry.

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Because symmetry is so universal, Group Theory is highly ubiquitous: it arises naturally not only in many fundamental areas of mathematics (like Geometry, Topology, Number Theory, Harmonic Analysis and more); but also in other areas of human study (like Virology, Chemistry, Physics, Computer Science, Cryptography..).

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The Extension Problem

Classify all of the finite groups.

In particular, we know the finite groups of order up as far as $2047 = 2^{11} - 1$. Here are the number of groups of order 2^k , for $k \le 10$.

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We will come back to this heuristic later in the talk, but for now...

Moral of the story: In general, the Extension Problem is intractable! For this reason, group theorists focus on other questions/areas, which aim to get us as close to a solution to the Extension Problem as possible. For the remainder of the talk, I will speak about three of these questions/areas:

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- 1. The Classification of Finite Simple Groups;
- 2. Asymptotic group theory;
- 3. Burnside's problems.

1. The Classification of Finite Simple Groups

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But as we saw before, we can classify some classes of finite groups (e.g. those of prime order; those of order p^2 or $2p^2$ for a prime p, and much more.)

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But why are the finite simple groups so important?

The Jordan–Hölder theorem

Theorem (Jordan–Hölder theorem)

Every finite group G has a composition series, i.e. a series

$$\{1_G\} = G_0 \trianglelefteq G_1 \trianglelefteq \cdots \trianglelefteq G_r = G$$

such that for each $1 \le i \le r$, the group G_i/G_{i-1} is simple. Moreover, although there can be different composition series, the length r, and the isomorphism classes of the factors G_i/G_{i-1} do not change. Thus, the multiset $\{\{G_i/G_{i-1} : 1 \le i \le r\}\}$ is well-defined, and is called the set of composition factors for G.

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E.g. The alternating groups A_n are simple for $n \ge 5$, while the cyclic groups of prime order are simple. The composition factors of S_n for $n \ge 5$ are $\{\{A_n, C_2\}\}$.

E.g. For p an odd prime, the composition factors of the dihedral group D_{2p} of order 2p are $\{\{C_p, C_2\}\}$.

For this reason, understanding the finite simple groups is crucial if one wants to get anywhere near the Extension Problem..

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For example, in Chapter 3 of MA3K4, we use Sylow's theorems to show that various groups (for example groups of order $4p^n$) <u>cannot</u> be simple, while in MA442, Sylow's theorems (and various other ideas, such as Burnside's transfer theorem) are used to classify the finite simple groups of order at most 500).

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Theorem (Burnside's $p^a q^b$ theorem)

Let p and q be primes, and let G be a group of order p^aq^b . Then G is not simple.

Then in the 1970s, there was an even more breathtaking advancement in our understanding of the finite simple groups.

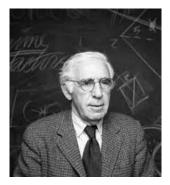
Theorem (The Odd Order Theorem; Feit & Thompson, 1970) Let G be a finite group of odd order. Then G is not simple.

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Theorem (The Odd Order Theorem; Feit & Thompson, 1970) Let G be a finite group of odd order. Then G is not simple.

Encouraged by this, together with some other deep results on finite simple groups like the Brauer-Fowler theorem, Daniel Gorenstein announced in 1972 an ambitious 16 step programme that would <u>classify</u> the finite simple groups.



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Side note: The last category is a set of finite simple groups of order between 7920 (the order of the *Mathieu group* M_{11}) and

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These groups do not fit naturally into any of the preceding three infinite families.

Gorenstein's programme

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In pursuit of a contradiction, one then chooses a finite simple group G such that G is not one of those listed above, and |G| is as small as possible.

From the previous theorems, we know that G has even order, and that |G| has at least 3 prime divisors. The Brauer-Fowler theorem also gives us information about the centralisers of elements of order 2 in G; while a method due to Bender gives certain restrictions on the maximal subgroups of G.

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Gorenstein's programme

In pursuit of a contradiction, one then chooses a finite simple group G such that G is not one of those listed above, and |G| is as small as possible.

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The minimality of |G| also shows that all subgroups of G have their composition factors lying in our know list of finite simple groups.

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<u>Also</u>, the finite simple groups on our list are complicated! We still don't know everything we need to know about them (for example, we still haven't been able to classify all of their maximal subgroups). This brings us nicely to..

Ah, a quick aside just before we move on..

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The final proof will be more uniform in approach, and will be shorter. (The plan is a volume of 11 books, comprising about 3000 pages in total. Book 9 is almost ready..).

2. Asymptotic group theory

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We've see a little more of this later on when we look at the Restricted Burnside Problem. There, group theorists don't try to *classify* the d-generated finite groups with exponent n. They just tried to count them.

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<i>G</i>	2	2 ²	2 ³	24	2 ⁵	2 ⁶	27	2 ⁸	2 ⁹	2 ¹⁰
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Conjecture (Erdös, 1965)

Let f(n) be the number of isomorphism classes of finite groups of order n. If $n, x \in \mathbb{N}$ with $n \leq 2^x$, then $f(n) \leq f(2^x)$.

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Both of these conjectures are still open.

The number of isomorphism classes of finite groups of order at most n

For a positive integer n with prime factorisation

 $n = p_1^{a_1} p_2^{a_2} \dots p_k^{a_k}$ (p_i distinct primes)

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For a functions g(n) and h(n), the notation h(n) = o(g(n)) means that $h(n)/g(n) \to 0$ as $n \to \infty$. For example, $n^{5/3} = o(n^2)$.

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In 1990, Pyber proved the following incredible result.

Theorem (Pyber, 1990)

Let $n \in \mathbb{N}$ and let $\mu := \mu(n)$. We have

$$f^*(n) \le n^{\mu^2/27+h(\mu)}$$

where $h(\mu) = o(\mu^2)$.

By work of Higman and Sims from the 1960s, the bound in the above theorem is "best possible".

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Let $n \in \mathbb{N}$. The number of subgroups of the symmetric group S_n is at most $2^{0.71n^2+h(n)}$, where $h(n) = o(n^2)$.

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He conjectured, however, that much more is true. (The following would be best possible).

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Theorem (Roney-Dougal & T., 2023)

Pyber's conjecture holds. In fact, the number of subgroups of the symmetric group S_n is at most $2^{n^2/16+cn^{3/2}}$, where c is some absolute constant.

Recall the probabilistic conjectures of Erdös and Pyber from earlier.

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An analogue for permutation groups was proposed by Kantor.

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Let $|\operatorname{Sub}(S_n)|$ and $|\operatorname{Sub}_2(S_n)|$ be the number of subgroups and 2-subgroups of S_n , respectively. Then $|\operatorname{Sub}_2(S_n)|/|\operatorname{Sub}(S_n)| \to 1$ as $n \to \infty$. That is, a random subgroup of S_n has order a power of 2.

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Theorem (Roney-Dougal & T., 2024)

Kantor's conjecture is not true. Indeed, there exists an absolute constant $\epsilon > 0$ such that $|\operatorname{Sub}_2(S_n)| / |\operatorname{Sub}(S_n)| < 1 - \epsilon$ for all $n \in \mathbb{N}$.

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One can take $\epsilon = 1/2^{16^2}$.

A number of questions arise from this theorem:

Question 1

Can we use this to make progress on the Erdős and Pyber conjectures?

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The final part of Question 2 is not "just for the sake of it". There is very practical motivation..

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The Group Isomorphism Problem (version we usually study)

Can we come up with an algorithm such that, given the multiplication tables of two groups G_1 and G_2 of order n as input, a computer can decide in a time which is polynomial in n, whether or not $G_1 \cong G_2$?

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In total, this takes time $n^{\log n + o(\log n)}$...

And we still can't do much better! The best general result to date is to due to Rosenbaum (2013), who showed that one can solve Grpl in time $n^{0.5 \log n + o(\log n)}$.

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An easy method

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In total, this takes time $n^{\log n + o(\log n)}$...

And we still can't do much better! The best general result to date is to due to Rosenbaum (2013), who showed that one can solve Grpl in time $n^{0.5 \log n + o(\log n)}$.

What has this got to with counting subgroups of finite groups?!

An intriguing new approach due to Gowers shows that an improved understanding of subgroup enumeration could lead to remarkable progress in Grpl.

The idea is as follows: For finite groups G and X, let $\operatorname{Sub}_{\cong X}(G)$ be the set of subgroups H of G with $H \cong X$. For $n \in \mathbb{N}$, let k(n) be the smallest positive integer such that for any two finite groups G_1 and G_2 of order n, we have $G_1 \cong G_2$ if and only if $|\operatorname{Sub}_{\cong X}(G_1)| = |\operatorname{Sub}_{\cong X}(G_2)|$ for all k(n)-generated finite groups X. Then the Grpl can be solved in time $n^{k(n)}$.

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Gowers' k(n) problem

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All we know at the moment is that $k(n) \leq \log n$.

3. Burnside's problems

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In mathematical terms, Group Theory is quite a young subject (Galois first defined a group in 1830, though Euler and Lagrange had already done some work on groups, under a different name, in the 18th century).

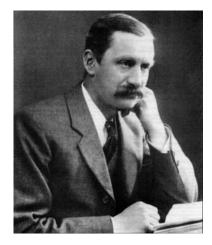
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By the late 1800s, we knew very little about abstract group theory. One of the most influential early group theorists was William Burnside.

In 1901, Burnside asked the following famous question.



Burnside's problem

Let G be a group which can be generated by d elements (G is said to be d-generated in this case), and such that $g^n = 1_G$ for all $g \in G$ (G is said to have exponent n in this case). Is G necessarily finite?

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Definition

Let G be a group, and let A be a non-empty subset of G. The subgroup of G generated by A, written $\langle A \rangle$, is defined to be

$$\langle A \rangle := \{a_1^{\epsilon_1} \dots a_m^{\epsilon_m} : m \in \mathbb{N}, a_i \in A, \epsilon_i \in \mathbb{Z}\}.$$

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So Burnside's problems says that if G is a group, there exists $A \subseteq G$ of size d such that every element can be written as a product of powers of elements of A, and every element of G has order at most n, then is G finite? Or, more informally, if G "operationally" finite, then is G finite?

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Thus, G is abelian. Hence, writing $A = \{x_1, \ldots, x_d\}$, we have

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Example 2 Suppose that G, d and n are as in Burnside's question, and n = 3 (i.e. G is a d-generated group in which $g^3 = 1_G$ for all $g \in G$).

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The same approach won't work (there are examples of finite groups of exponent 3 which are <u>not</u> abelian). But is there anything we can do?

Yes! Claim: Let G be a group of exponent 3. Then G is "almost abelian". More precisely, a commutes with $a^b := b^{-1}ab$ for all $a, b \in G$.

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Proof of claim: Let G be as in the statement, and let a and b be elements of G. Then we have

$$1_G = (ba)^3 = bababa$$

= $b(b^2b^{-2})a(b^2b^{-2})baba$
= $b^3(b^{-2}ab^2)(b^{-2}b)aba$
= $(bab^{-1})(b^{-1}ab)a$ since $b^3 = 1_G$, and hence $b = b^{-2}$.
= $(bab^{-1})(b^2ab^{-2})a$.

We therefore have $a^{-1} = a^{b^{-1}}a^b$. Replacing *b* by b^{-1} gives $a^{-1} = a^b a^{b^{-1}}$. Thus,

$$a^{b^{-1}}a^b = a^b a^{b^{-1}}$$
 for all $a, b \in G$.

So raising both sides to the b^2 yields:

$$a^b a = a a^b$$
 for all $a, b \in G$.

$$N := \langle a_d^b : b \in G \rangle$$

is abelian, hence isomorphic to C_3^m for some m (by the same argument as used in Example 1). Since $N \subseteq G$ and G/N can be generated by $\langle a_1N, \ldots, a_{d-1}N \rangle$, an easy inductive argument shows that G is finite.

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So in summary, recalling:

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- ▶ What about the case *n* = 5?

NOBODY KNOWS ..

Burnside's problem in exponent 5 is notoriously difficult: we know almost nothing about the problem in this case.

Even if we restrict to the case d = 2, we still have no idea what happens.

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Theorem (Adian and Novikov, 1968)

For every odd number n with n > 4381, there exist infinite, finitely generated groups of exponent n. Thus, the answer to Burnside's problem is NO.

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These are called *Tarski Monsters*, and they exist for every prime $p > 10^{75}$.

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Because of this difficulty, and in general the lack of progress made in the first 30 years after the statement of Burnside's problem, mathematicians started to ask a weaker question in the 1930s. This became known as the *Restricted Burnside problem*:

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The Restricted Burnside problem

With the Hall-Higman reduction in mind, group theorists began working furiously on the restricted Burnside problem in prime power exponent. This allows one to assume that the group G one is working in is a finite p-group, for some prime p (i.e. G has p-power order).

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The first significant breakthrough came from Kostrikin in 1959.

Theorem (Kostrikin, 1959)

Fix $d \in \mathbb{N}$ and a prime p. There are only finitely many finite d-generated groups G of exponent p. That is, the Restricted Burnside problem has a positive solution for prime exponents.

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Kostrikin's idea used an intriguing connection between finite groups of *p*-power order, and *Lie algebras* (on which Warwick's Adam Thomas is one of the world's foremost experts!)

Definition

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- (c) [v, v] = 0 for all $v \in L$.
- (d) The Jacobi identity [u, [v, w]] + [w, [u, v]] + [v, [w, u]] = 0 holds for all u, v, w ∈ L.

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$$[v, v] = 0$$
 for all $v \in L$.

(d) The Jacobi identity [u, [v, w]] + [w, [u, v]] + [v, [w, u]] = 0 holds for all u, v, w ∈ L.

The standard example is $L := M_n(\mathbb{F})$, the set of $(n \times n)$ -matrices over \mathbb{F} , where [A, B] := AB - BA.

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Let \mathbb{F} be a field. A *Lie algebra over* \mathbb{F} is a pair $(L, [\cdot, \cdot])$ where

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But what has this got to do with finite *p*-groups?!

Let G be a finite group of exponent p (i.e. $g^p = 1_G$ for all $g \in G$). The commutator subgroup $\gamma_2(G) := [G, G]$ is defined by

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We now define $\gamma_1(G) := G$, and for $i \ge 2$,

$$\gamma_i(G) := \langle [x,g] : x \in \gamma_{i-1}(G), g \in G \rangle.$$

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We then define $L(G) := \gamma_1(G)/\gamma_2(G) \oplus \gamma_2(G)/\gamma_3(G) \oplus \ldots \oplus \gamma_n(G)/\gamma_{c+1}(G); \text{and we set}$ $[v\gamma_i(G), w\gamma_j(G)] := [v, w]\gamma_{i+j \pmod{n}}(G).$

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The case of general prime power exponent $d := p^m$ was done by Zelmanov in the late 1980s. He used a similar connection to the above, but needed to work with *Lie rings* rather than Lie algebras, as $\mathbb{Z}/p^m\mathbb{Z}$ is a field if and only if m = 1.

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In 1994, Zelmanov was awarded the Fields medal for his work.



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In my view, another huge aspect of its beauty is how quickly one can get from the definition of a group, to deep and important problems.

For example, if you can prove that a 2-generated group G in which $g^5 = 1_G$ for all $g \in G$, is finite, then you will (without a doubt) win a Fields medal...